Fortran at the Intersection:

Synergies Arising from the Interplay Between Paradigms

Damian Rouson







Outline

- Introduction
 - Case Studies
 - Intersectionality
 - Acknowledgments

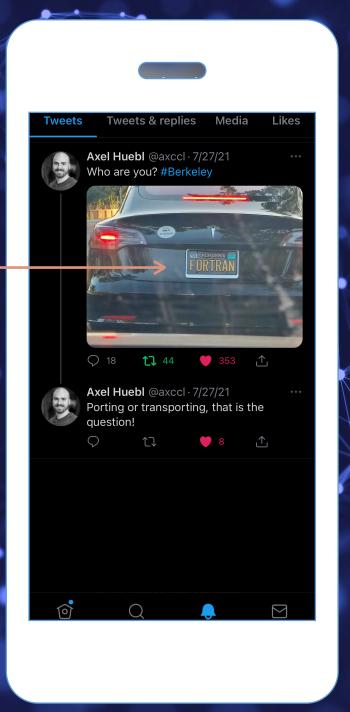


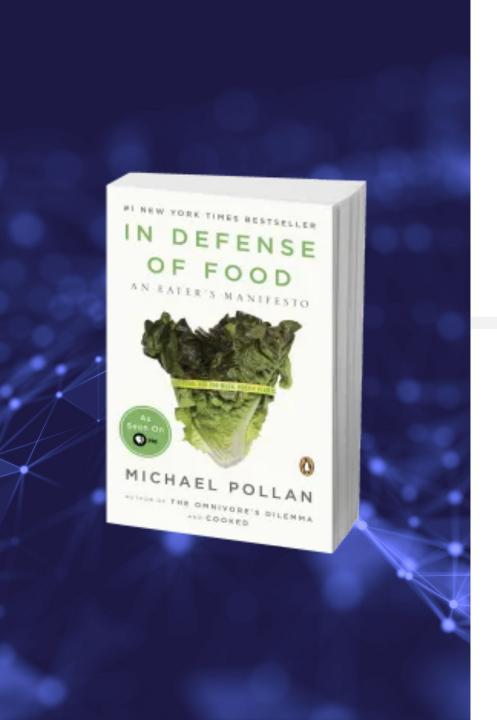
Outline

- Introduction
 - ** Background: Fortran at the Intersection.
 - Motivation: A preliminary summary.
 - Case Studies
 - Intersectionality
 - Acknowledgments









Eat food.

Not too much.

Mostly plants.



In Defense of Software: A Developer's Manifesto

Write software.

Not too much.

Mostly pure functions.

Outline

- Introduction
 - Case Studies: Open-Source Software
 - **Rocket Science**
 - Synergies at the Intersection of Paradigms
 - # FEATS, Dag, and Assert
 - Intersectionality
 - Acknowledgments



Rocket Science

An educational mini-app for predicting solid rocket motor thrust.

Instructions to client:

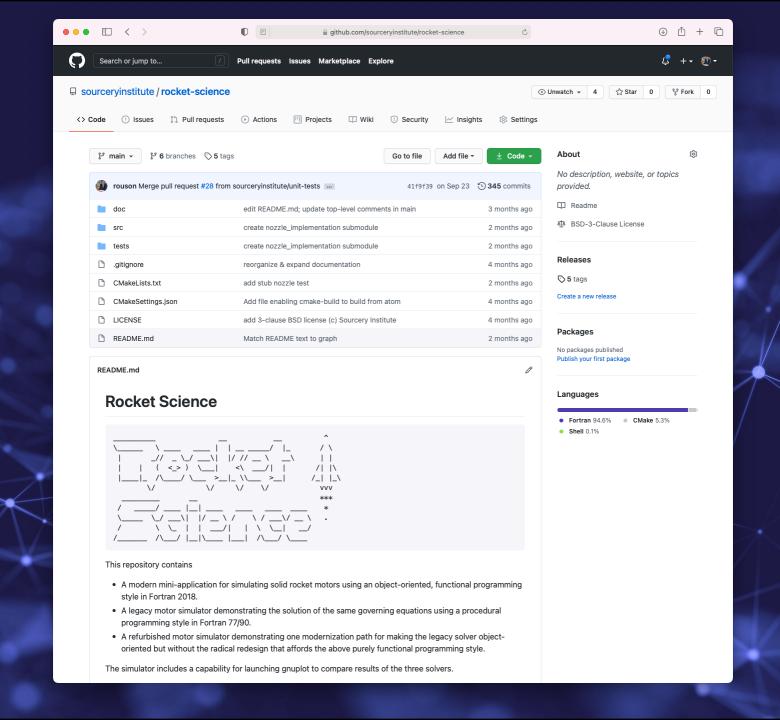
- "Write a small proxy application capturing an algorithm that plays an important role in your production application."
- "Write the proxy app exactly as you would the day before you met us."

Result:

♣ A procedural Fortran 90 code in which subroutines with 0-1 arguments collectively use one common module containing approximately 40 public module variables.

Goals:

- ▶ To demonstrate an evolutionary approach for refactoring the proxy app using modern programming paradigms.
- ▶To use the process as a case study in a training course.



Programming Paradigms

Functional Programming

Object-Oriented Programming

Parallel Programming Programming by contract

What is a pattern in one language might be a feature in another language, including a future version of the same language.

Thus, patterns might be strong candidates for future inclusion in the language: some aspects of the current considerations for generic programming resemble programming by contract.

A Functional Programming Pattern

The Pattern:

- Write pure, including elemental, functions.
- ♣ Associate names with expressions, including function references.
- ★ Construct whole objects with generic interfaces serving as user-defined structure constructors → no setters.
- Write expressions based on pure userdefined operators.

```
src - vim main.f90 - 135×33
    !! A mini-application for rocket motor simulation:
     !! demonstrating an object-oriented, functional programming style in Fortran 2018
       with automatic graphing of results in gnuplot for comparison against a legacy,
     !! procedural simulator written in Fortran 90.
     use motor_interface, only : motor_t
     use state_interface, only : state_t
     use kind_parameters, only : rkind
    implicit none
    real(rkind), parameter :: zero=0._rkind
     character(len=*), parameter :: input_file="rocket.inp"
12
     type(motor t) motor
    type(state_t) state !! state variables updated at each time step: mass, energy, time, and burn depth
    type(state_t), allocatable :: history(:)
16
     call motor%define(input_file)
17
18
     associate(chamber => motor%chamber())
       associate(gas => chamber%gas())
        call state%define(input_file, R_gas=gas%R_gas(), c_v=gas%c_v(), volume=chamber%initial_volume(), time=zero, burn_depth=zero)
       end associate
    end associate
    associate(dt => motor%dt() )
      history = [state]
       do while(state%time() < motor%t_max())</pre>
        state = state + motor%d dt(state)*dt
        history = [history, state]
       end do
    end associate
                                                                                                                                    Top
```

A Functional Programming Pattern

The Benefits:

- ▶ Clarity: the reader sees inputs map to outputs without reading the interface.
- Robustness*: associating with expressions provides immutable state with clear, limited scoping.
- Performance: only pure procedures may be invoked inside do concurrent blocks, easier to write asynchronous parallel code, + other potential optimizations.
- * This reduces or eliminates whole categories of errors that arise form the traditional declare-and-define pattern, e.g., no more mistaken use of uninitialized data or mistakenly overwriting data.

```
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     use kind_parameters, only : rkind
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     character(len=*), parameter :: input_file="rocket.inp"
12
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    type(state_t) state !! state variables updated at each time step: mass, energy, time, and burn depth
     type(state_t), allocatable :: history(:)
16
     call motor%define(input_file)
18
     associate(chamber => motor%chamber())
       associate(gas => chamber%gas())
        call state%define(input_file, R_gas=gas%R_gas(), c_v=gas%c_v(), volume=chamber%initial_volume(), time=zero, burn_depth=zero)
       end associate
23
     end associate
24
     associate(dt => motor%dt() )
      history = [state]
       do while(state%time() < motor%t max())</pre>
         state = state + motor%d dt(state)*dt
        history = [history, state]
30
       end do
    end associate
```

A Functional Programming Pattern

The Roadblocks:

- ▶ Functions can only have one result.
- ▶ Producing output for debugging is cumbersome.
- ▶ If copies replace mutations, memory could be used inefficiently.

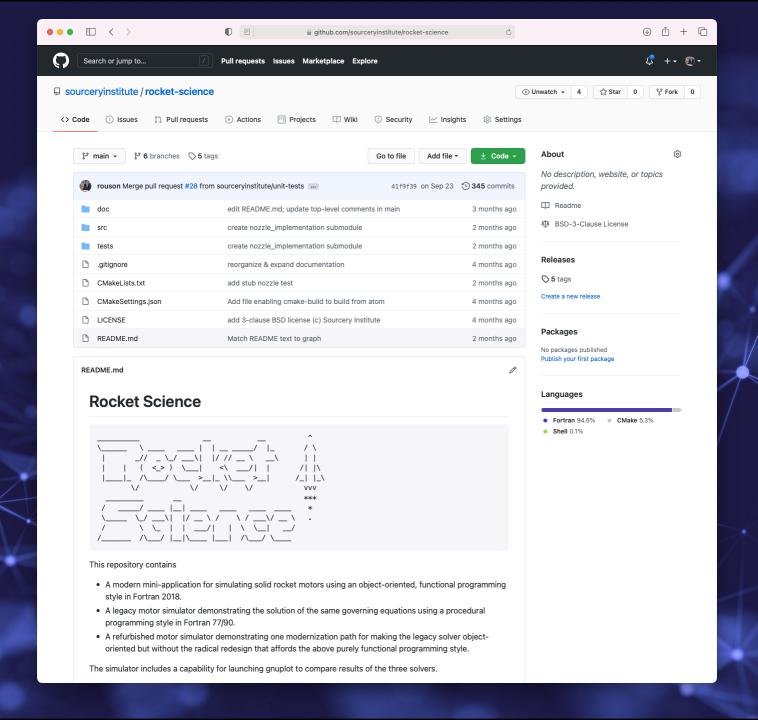
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src - vim main.f90 - 135×33
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       end associate
     end associate
     associate(dt => motor%dt() )
       history = [state]
       do while(state%time() < motor%t_max())</pre>
         state = state + motor%d dt(state)*dt
        history = [history, state]
       end do
    end associate
                                                                                                                                    Top
```

Rocket Science

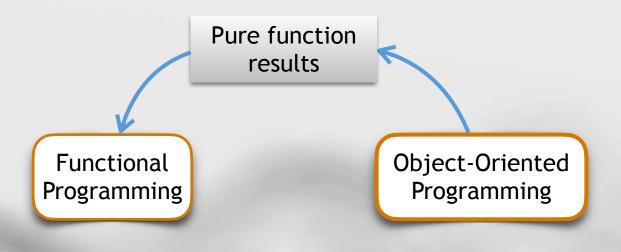
An educational mini-app for predicting solid rocket motor thrust.

Steps to go from software archaeology to modernity:

- Set up a git repository.
- Add an fpm manifest.
- **Establish** a reference implementation:
- Set up a test suite using Vegetables.
- ▶ Define GitHub Actions for
 - Automating continuous integration testing.
 - Generating ford documentation.
 - Deploying documentation to GitHub Pages.
- ▶ Convert module variables to dummy arguments with intent.
- Separate procedure interface bodies (in modules) from procedure definitions (in submodules).
- Convert subroutines to pure functions:
 - intent(in): no change.
 - intent(out): encapsulate in a derived type result.
 - intent(inout) -> intent(in) + result component.



Synergies at the Intersection



Parallel Programming

Programming by contract

Patterns resolve tensions in program design:

Using objects as function results resolves the tension between needing multiple results but only being allowed one.

This approach often also leads to a natural encapsulation strategy: if one procedure produces multiple values, it's likely those values are related in a way that makes a reasonable choice of abstractions.

FEATS

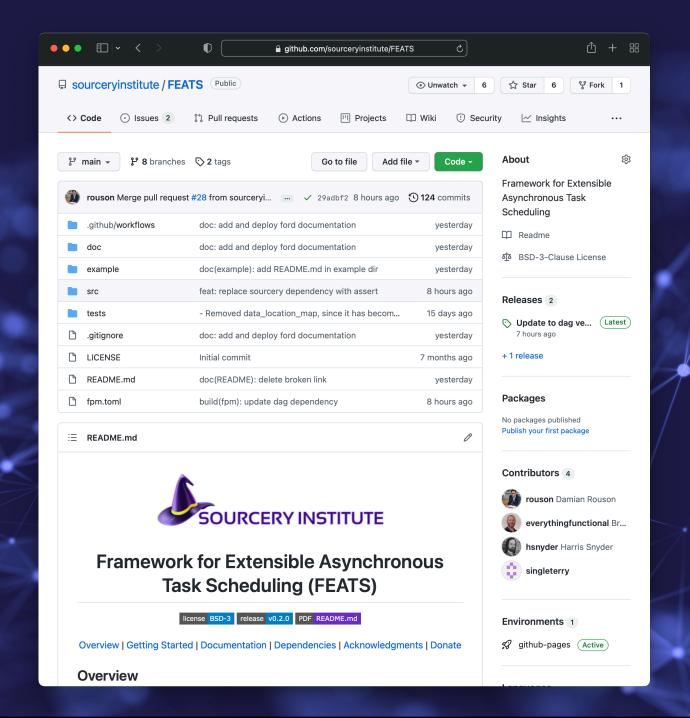
Framework for Extensible Asynchronous Task Scheduling

Execution:

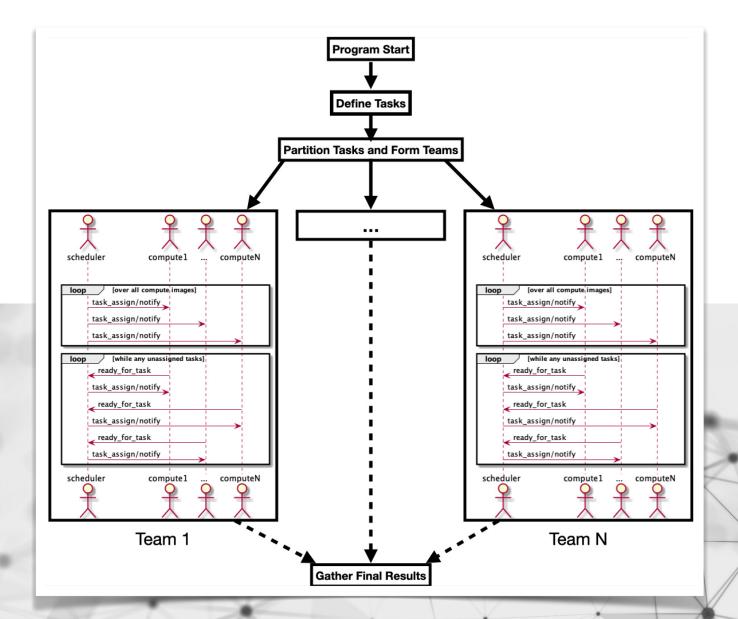
- In each team, establish one scheduler image and one or more compute images.
- Schedulers post task_assigned events to compute images in an order that respects dependencies in a directed acyclic graph (DAG).
- Compute images post ready_for_next_task events to scheduler.
- A task_payload_map_t abstraction maps task task identifiers to locations in a payload_t mailbox coarray.

Initial target applications:

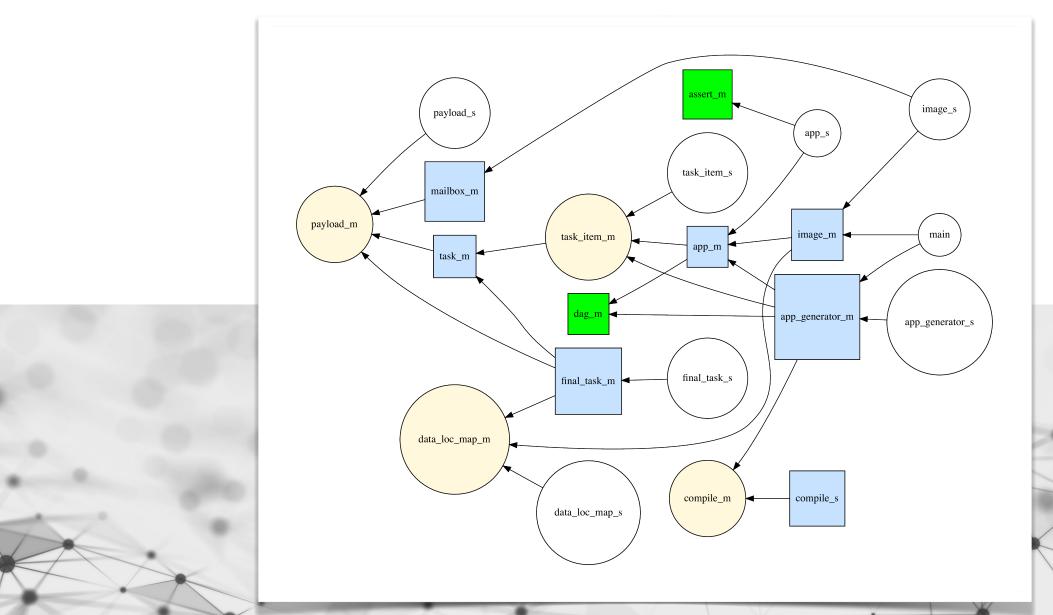
- NASA's Online Tool for the Assessment of Radiation in Space (OLTARIS)
- ▶ NCAR's Intermediate Complexity Atmospheric Research (ICAR) model: work-sharing/work-stealing.
- ▶ Fortran Package Manager: parallel builds.



UML Sequence Diagram



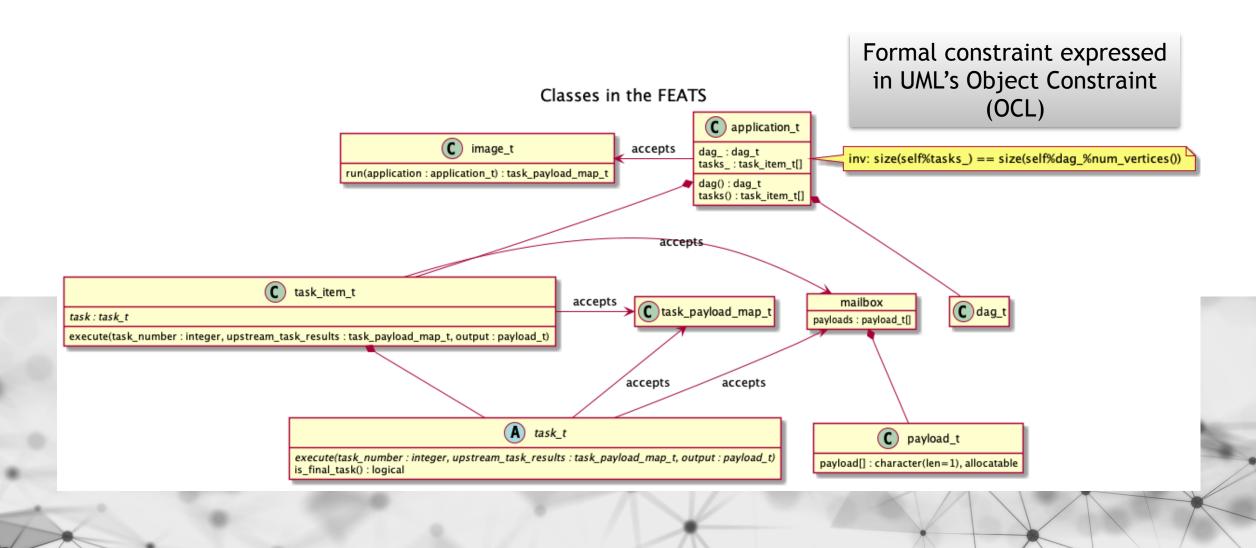
FEATS Module Dependency DAG



FEATS Parallel Traversal of the FEATS DAG

```
[rouson:~/Repositories/sourceryinstitute/feats] main+* 39s 130 ±
```

UML Class Diagram with a Simple Contract



Application_t Interface & Implementation

```
application_s.f90 — ~/Repositories/sourceryinstitute/feats
         application_m.f90
                                                                                      application_s.f90
       type application t
                                                                                 submodule(application_m) application_s
                                                                                   use assert m, only : assert
         private
                                                                                   implicit none
         type(dag_t) dag_ !! Describes the dependencies between tasks
         type(task_item_t), allocatable :: tasks_(:) !! tasks to be exect 5
         private
                                                                                   module procedure construct
         procedure, public :: dag
                                                                                     call assert(size(tasks)==dag%num vertices(), &
         procedure, public :: tasks
                                                                                       "application(construct): size(tasks)==dag%num_vertices()")
       end type
                                                                                     application%dag = dag
                                                                                     application%tasks_ = tasks
       interface application t
                                                                                   end procedure
         pure module function construct(dag, tasks) result(application)
                                                                                   module procedure dag
           implicit none
                                                                                     dag = self%dag_
           type(dag t), intent(in) :: dag
                                                                                   end procedure
           type(task_item_t), intent(in) :: tasks(:)
           type(application t) application
                                                                                   module procedure tasks
         end function
                                                                                     tasks = self%tasks_
                                                                                   end procedure
       end interface
                                                                                 end submodule application s
       interface
         pure module function dag(self)
           implicit none
           class(application_t), intent(in) :: self
           type(dag_t) :: dag
         end function
         pure module function tasks(self)
           implicit none
           class(application_t), intent(in) :: self
           type(task_item_t), allocatable :: tasks(:)
         end function
                                                                                             • LF UTF-8 Fortran - Free Form & main & Fetch  GitHub - Git (2)
fpm build src/application_s.f90* 0 11 ▲ 0 ⊙ 0 9:7
```

Application Construction

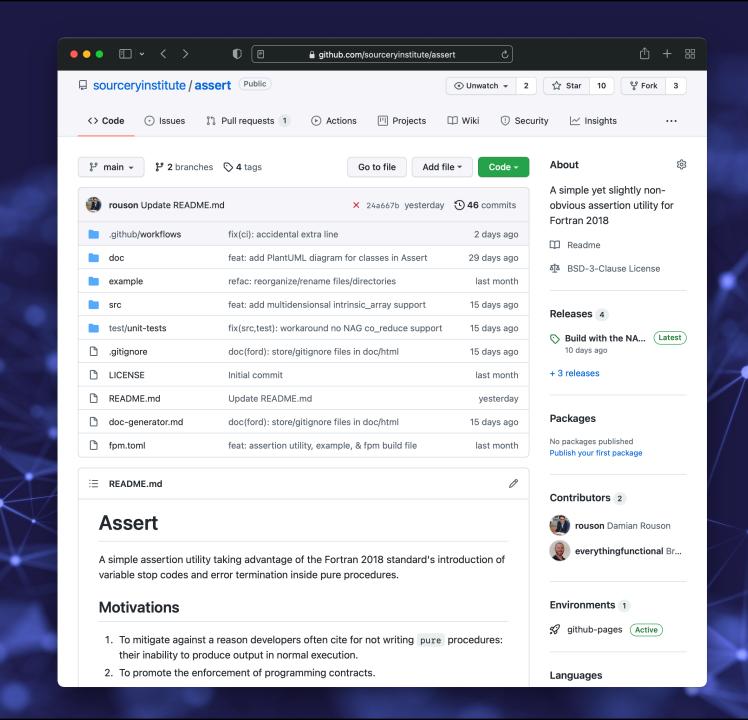
```
• • •
                                                  application generator m.f90 - ~/Repositories/sourcervinstitute/feats
     application generator m.f90
                       character(len=*),
                                                    parameter :: external_ = 'shape=square,fillcolor="green",style=filled'
                                                                           = 'shape=circle.fillcolor="white".style=filled'
                       character(len=*),
                                                    parameter :: root
                       character(len=*).
                                                                           = 'shape=square,fillcolor="SlateGray1",style=filled'
                                                    parameter :: branch
                       character(len=len(branch)), parameter :: leaf
                                                                           = 'shape=circle.fillcolor="cornsilk".style=filled'
                       type(dag_t) feats
                       type(varying_string) name_string(size(names))
                       type(task_item_t), allocatable :: tasks(:)
                       name string = var_str(names)
                       feats = &
                           dag_t([ &
                                 vertex_t([integer::],
                                                                                                                 var_str(external_)) &
                                , vertex_t([integer:: ],
                                                                                   name string(dag m),
                                                                                                                 var_str(external_)) &
                                , vertex_t([integer::],
                                                                                   name_string(payload_m),
                                                                                                                                    ) &
                                                                                                                                    ) &
                                                                                   name_string(compile_m),
                                                                                                                 var_str(leaf)
                                                                                                                                    3 (
                                , vertex_t([integer::],
                                                                                   name_string(data_loc_map_m), var_str(leaf)
                                , vertex_t([payload_m],
                                                                                                                                   ) &
                                                                                   name_string(task_m),
                                                                                                                 var_str(branch)
                                                                                                                                    . &
                                , vertex_t([task_m],
                                                                                   name string(task item m),
                                , vertex_t([dag_m, task_item_m],
                                                                                   name_string(app m),
                                                                                                                                   ) &
                                                                                                                 var_str(branch)
                                , vertex t([app m, dag m, task item m, compile m], name string(app generator m), var str(branch)
                                , vertex t([app m, data loc map m],
                                                                                    name_string(image_m),
                                                                                                                 var_str(branch)
                                , vertex_t([app_generator_m, image_m],
                                                                                                                                    ) &
                                                                                                                                    ) &
                                , vertex_t([task_item_m],
                                                                                   name string(task item s),
                                , vertex_t([compile_m],
                                                                                                                                    ) &
                                                                                                                 var_str(branch)
                                                                                                                                    ) &
                                , vertex_t([app_generator_m],
                                                                                   name_string(app_generator_s), var_str(root)
                                                                                                                                    3 (
                                , vertex_t([data_loc_map_m],
                                                                                   name_string(data_loc_map_s), var_str(root)
                                , vertex_t([payload_m],
                                                                                   name_string(payload_s),
                                                                                                                                    ) &
                                , vertex_t([app_m, assert_m],
                                                                                   name string(app s),
                                                                                                                                    . &
                                , vertex_t([payload_m],
                                                                                                                 var_str(branch)
                                                                                                                                    ) &
                                                                                                                                    ) &
                                , vertex_t([image_m, mailbox_m],
                                                                                   name_string(image_s),
                                                                                                                                    ) &
                                , vertex_t([data_loc_map_m, payload_m, task_m],
                                                                                   name_string(final_task_m),
                                , vertex_t([final_task_m],
                                                                                                                                    ) &
                       tasks = [(task_item_t(compile_task_t(name_string(i))), i = 1, size(names))]
                   end block
 • LF UTF-8 Fortran - Free Form & main & Fetch  GitHub  Git (5)
```

Assert

A simple yet slightly nonobvious assertion utility for Fortran 2018

Goals:

- ♣ To promote contract enforcement via runtime constraint checking.



Assert

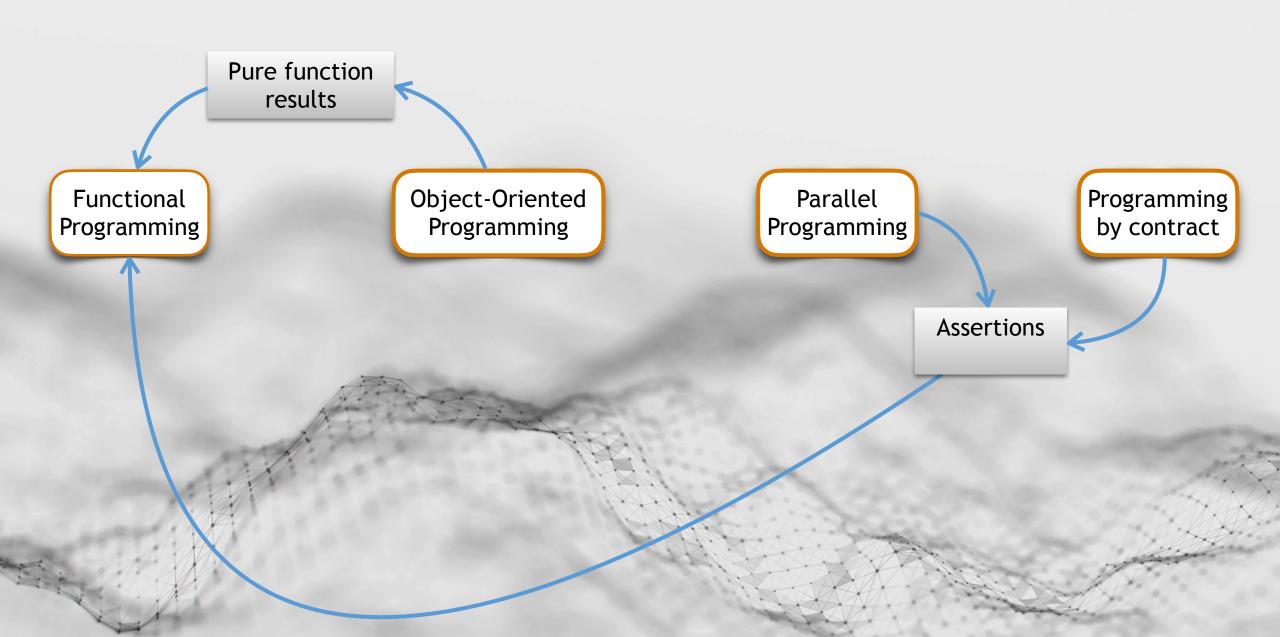
A simple yet slightly nonobvious assertion utility for Fortran 2018

Approach:

- ▶ Express constraints in logical expression as actual arguments to a pure assert subroutine.
- ♠ Execute error stop upon assertion failure.
- Bracket all executable statements in in the assert() subroutine by a conditional evaluation of a parameter logical ASSERTIONS variable
- -> facilitate dead-code removal optimization

```
assert_m.F90 — ~/Repositories/sourceryinstitute/asser
     assert m.F90
#ifndef USE_ASSERTIONS
# define USE_ASSERTIONS .true.
#endif
  logical, parameter :: enforce_assertions=USE_ASSERTIONS
 interface
    pure module subroutine assert(assertion, description, diagnostic data)
      implicit none
      logical, intent(in) :: assertion
      character(len=*), intent(in) :: description
      class(*), intent(in), optional :: diagnostic data
    end subroutine
  end interface
end module
  src/assert_m.F90 🔞 2 🛦 0 🕫 0 24:1 • LF 🛕 2 deprecations UTF-8 Fortran - Free Form 👂 rm-ford-6.0.0-workaround 🧲 Fetch 🕡 GitHub 🕟
```

Synergies at the Intersection



Assert

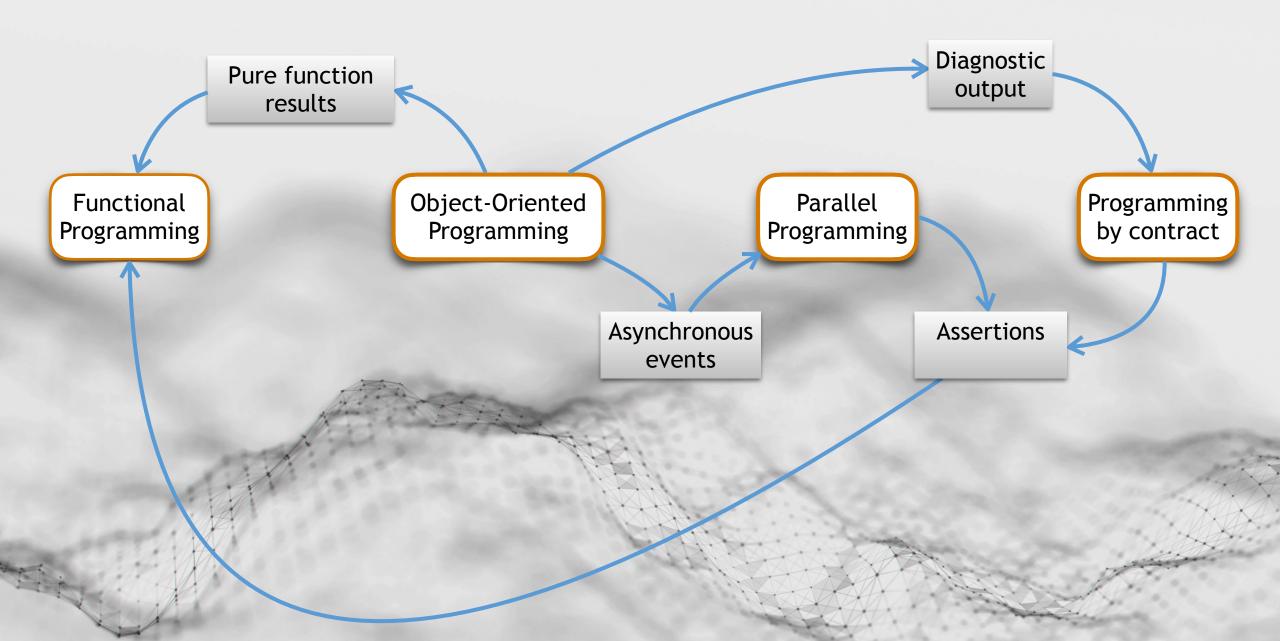
A simple yet slightly nonobvious assertion utility for Fortran 2018

Approach:

- ▶ Express constraints in logical expression as actual arguments to a pure assert subroutine.
- Execute error stop upon assertion failure.
- Bracket all executable statements in in the assert() subroutine by a conditional evaluation of a parameter logical ASSERTIONS variable
- -> facilitate dead-code removal optimization
- ♣ Optionally output diagnostic data as scalar default character expression stop codes.

```
assert_s.f90 — ~/Repositories/sourceryinstitute/assert
           assert s.f90
       module procedure assert
         use characterizable m, only : characterizable t
         character(len=:), allocatable :: header, trailer
         toggle_assertions: &
         if (enforce_assertions) then
           check assertion: &
           if (.not. assertion) then
             associate(me=>this_image()) ! work around gfortran bug
              header = 'Assertion "' // description // '" failed on image ' // string(me)
             end associate
             represent_diagnostics_as_string: &
             if (.not. present(diagnostic_data)) then
              trailer = "(none provided)"
              select type(diagnostic data)
                type is(character(len=*))
                type is(complex)
                  trailer = string(diagnostic data)
                type is(integer)
                  trailer = string(diagnostic data)
                type is(logical)
                  trailer = string(diagnostic_data)
                type is(real)
                  trailer = string(diagnostic data)
                class is(characterizable_t)
                  trailer = diagnostic data%as character()
```

Synergies at the Intersection



Abstract Calculus Pattern

Standard semantic requirement: operands must have the intent(in) attribute.

-> not merely syntactic sugar.

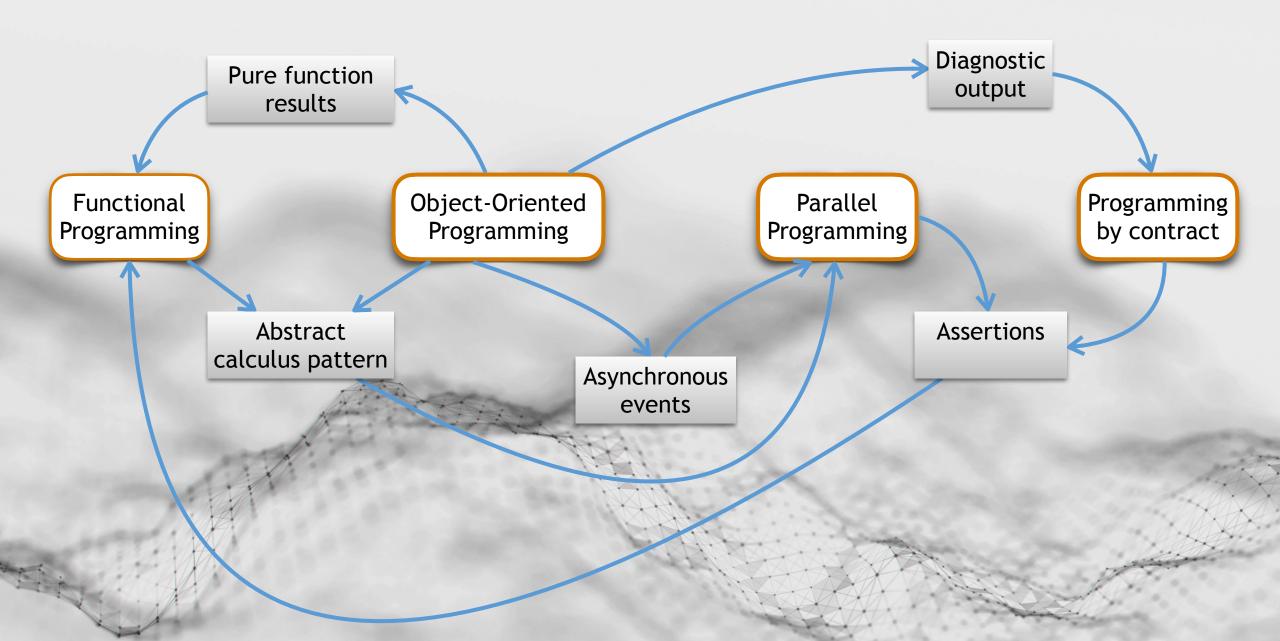
A nudge in the direction of functional programming?

Performance implication: fully asynchronous expression evaluation even with communication.

Image control happens in the one non-functional step: the assignment.

```
User-defined, purely functional operators
  (.grad.p)/rho + nu*(.laplacian.u) -(u.dot.(.grad.u))
             Distributed objects
```

Synergies at the Intersection



Outline

- Introduction
 - Case Studies
 - Intersectionality: Notes from an underdog
 - ****** A social science definition.
 - ** A diverse workforce does work differently...
 - And does different.





Prof. Kimberle Crenshaw **UCLA School of Law**

Intersectionality

University of Chicago Legal Forum Volume 1989 | Issue 1 Demarginalizing the Intersection of Race and Sex: A Black Feminist Critique of Antidiscrimination Doctrine, Feminist Theory and Antiracist Politics Follow this and additional works at: http://chicagounbound.uchicago.edu/uclf Crenshaw, Kimberle () "Demarginalizing the Intersection of Race and Sex: A Black Feminist Critique of Antidiscrimination Doctrine, Feminist Theory and Antiracist Politics," University of Chicago Legal Forum: Vol. 1989; Iss. 1, Article 8. Available at: http://chicagoumbound.arking.oc.du/u/f/v191899/iss.1) This Article is brought to you for free and open access by Chicago Unbound. It has been accepted for inclusion in University of Chicago Legal Forus

Google Scholar citation count: 21,620

Demarginalizing the Intersection of Race and Sex: A Black Feminist Critique of Antidiscrimination Doctrine, Feminist Theory and Antiracist Politics

Kimberle Crenshawt

One of the very few Black women's studies books is entitled All the Women Are White, All the Blacks Are Men, But Some of Us are Brave.\(^1\) I have chosen this title as a point of departure in my efforts to develop a Black feminist criticism\(^2\) because it sets forth a problematic consequence of the tendency to treat race and gender as mutually exclusive categories of experience and analysis.\(^3\) In this talk, I want to examine how this tendency is perpetuated by a single-axis framework that is dominant in antidiscrimination law and that is also reflected in feminist theory and antiracist politics.

I will center Black women in this analysis in order to contrast the multidimensionality of Black women's experience with the single-axis analysis that distorts these experiences. Not only will this juxtaposition reveal how Black women are theoretically erased, it will also illustrate how this framework imports its own theoretical limitations that undermine efforts to broaden feminist and an-

[†] Acting Professor of Law, University of California, Los Angeles Law School.

Gloria T. Hull, et al, eds (The Feminist Press, 1982).

⁸ For other work setting forth a Black feminist perspective on law, see Judy Scales-Trent. Black Women and the Constitution: Finding Our Place, Asserting Our Rights (Voices of Experience: New Responses to Gender Discourse), 24 Harv CR-CL L Rev 9 (1989); Regina Austin, Sapphire-Boundf, forthcoming in Wise Women's L J (1989); Angela Harris, Race and Essentialism in Feminist Legal Theory (unpublished manuscript on file with author); and Paulette M. Caldwell, A Hair Piece (unpublished manuscript on file with author).

³ The most common linguistic manifestation of this analytical dilemma is represented in the conventional usage of the term "Blacks and women." Although it may be true that some people mean to include Black women in either "Blacks" or "women," the context in which the term is used actually suggests that often Black women are not considered. See, for example, Elizabeth Spelman, The Inessential Woman 114-16 (Beacon Press, 1988) (discussing an article on Blacks and women in the military where "the racial identity of those identified as "women' does not become explicit until reference is made to Black women, at which point it also becomes clear that the category of women excludes Black women"). It seems that if Black women were explicitly included, the preferred term would be either "Blacks and white women" or "Black men and all women."



How do we tell truths that might hurt?

Edsger W.Dijkstra, 18 June 1975

from https://www.cs.utexas.edu/users/EWD/ewd04xx/EWD498.PDF

Sometimes we discover unpleasant truths. Whenever we do so, we are in difficulties: suppressing them is scientifically dishonest, so we must tell ther but telling them, however, will fire back on us. If the truths are sufficiently impalatable, our audience is psychically incapable of accepting them and w will be written off as totally unrealistic, hopelessly idealistic, dangerously revolutionary, foolishly gullible or what have you. (Besides that, telling such truths is a sure way of making oneself unpopular in many circles, and, as such, it is an act that, in general, is not without personal risks. Vide Galileo Galilei.....)

Computing Science seems to suffer severely from this conflict. On the wholit remains silent and tries to escape this conflict by shifting its attention. (Fo instance: with respect to COBOL you can really do only one of two things: fit the disease or pretend that it does not exist. Most Computer Science Departments have opted for the latter easy way out.) But, Brethern, I ask yo is this honest? Is not our prolonged silence fretting away Computing Sciencintellectual integrity? Are we decent by remaining silent? If not, how do we speak up?

To give you some idea of the scope of the problem I have listed a number o such truths. (Nearly all computing scientists I know well will agree without hesitation to nearly all of them. Yet we allow the world to behave as if we di not know them....)

- Programming is one of the most difficult branches of applied mathematics; the poorer mathematicians had better remain pure mathematicians.
- The easiest machine applications are the technical/scientific computations.
- The tools we use have a profound (and devious!) influence on our thinking habits, and, therefore, on our thinking abilities.
- FORTRAN --"the infantile disorder"--, by now nearly 20 years old, is hopelessly inadequate for whatever computer application you have in mind today: it is now too clumsy, too risky, and too expensive to use.

An Underdog





1975-2001

2005

2021

"Music lovers are at high risk of being inspired by this exploration of the connections between music and physics." -Wall Street Journal THE SECRET LINK BETWEEN MUSIC AND THE STRUCTURE OF THE UNIVERSE JAZZ PHYSICS

A diverse workforce does work differently...

"... Then, out of the blue, Salam says to Jim: 'One day, when your people do physics, it will be like jazz.'

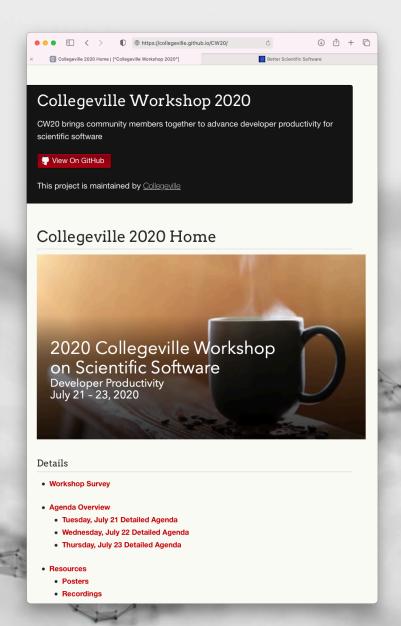
What a great compliment, an affirmation and acknowledgment of the improvisational, inclusive, cultural, and intellectual contributions of this music called jazz."

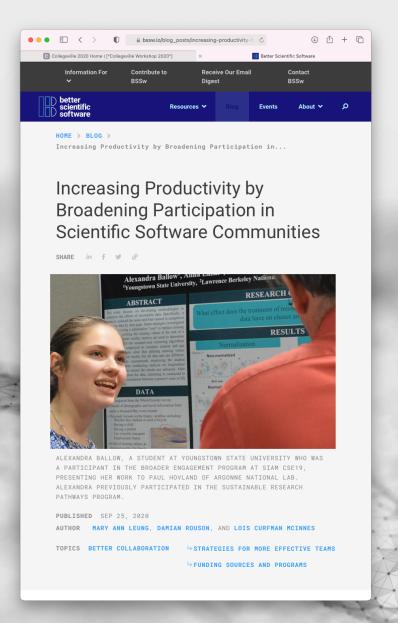
—Stephon Alexander (*The Jazz of Physics*, p. 231) quoting physics Nobel laureate Abdus Salam as relayed by physicist Jim Gates.

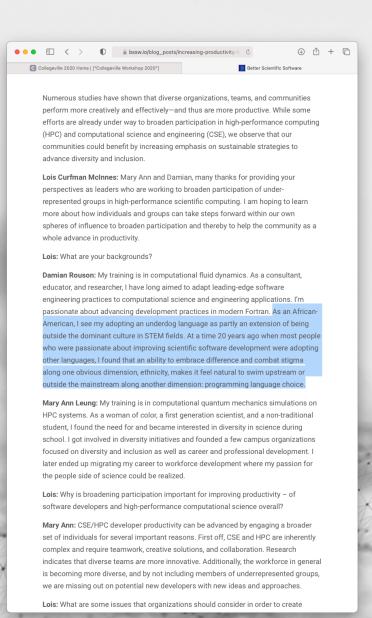
"It occurred to me by intuition, and music was the driving force behind that intuition. My discovery was the result of musical perception."

Albert Einstein (when asked about his theory of relativity)

... and does different work







Acknowledgements

- Zaak Beekman, ParaTools, Inc.
- ▶ Janne Blomquist, GCC
- ♣ Aurelien Bouteiller, University of Tennessee
- ▶ Tobias Burnus, Germany
- C. Carson, Stanford University
- Alessandro Fanfarillo, AMD
- Salvatore Filippone, University of Rome Tor Vergata
- Brian Friesen, Berkeley Lab
- ▶ Daniel Celis Garza, Oxford University
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